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## Tutorial 5

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**Problem 1.** (*basic requirements for cryptographic hash functions*) Using a block cipher  $E_K(x)$  with block length  $k$  and key  $K$ , a hash function  $h(m)$  is provided in the following way:

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Append  $m$  with zero bits until it is a multiple of  $k$ , divide  $m$  into  $n$  blocks of  $k$  bits each.
 $c \leftarrow E_{m_0}(m_0)$ 
for  $i$  in  $1..(n - 1)$  do
     $d \leftarrow E_{m_0}(m_i)$ 
     $c \leftarrow c \oplus d$ 
end for
 $h(m) \leftarrow c$ 

```

- Does this function fulfill the basic requirements for a cryptographic hash function?
- Can these requirements be fulfilled by replacing the operation XOR ( $\oplus$ ) by AND ( $\odot$ )?

**Problem 2.** (*proof of Example 10.2*) Complete the proof of Example 10.2 from the lecture notes. Show that from

$$k(x_1 - x'_1) \equiv x'_0 - x_0 \pmod{p-1}$$

the discrete logarithm  $k = \log_a(b) \pmod{p}$  can be efficiently computed.

**Problem 3.** (*Collision in hash functions*) Consider the following function:

$$h : \{0, 1\}^* \rightarrow \{0, 1\}^*, k \mapsto \left( \left[ 10000 \left( (k)_{10} (1 + \sqrt{5}) / 2 - \lfloor (k)_{10} (1 + \sqrt{5}) / 2 \rfloor \right) \right] \right)_2.$$

Here,  $\lfloor x \rfloor$  is the floor function of  $x$  (round down to the next integer smaller than  $x$ ). For computing  $h(k)$ , the bitstring  $k$  is identified with the positive integer it represents. The result is then converted to binary representation.

(example:  $k = 10011$ ,  $(k)_{10} = 19$ ,  $h(k) = (7426)_2 = 1110100000010$ )

- Determine the maximal length of the output of  $h$ .
- Give a collision for  $h$ .