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Problem 1. (decpiher Blum-Goldwasser) Bob receives the following cryptogram from Alice:

 $c = (101010111000011010001011100101111100110111000, x_{t+1} = 1306)$

The message m has been encrypted using the Blum-Goldwasser cryptosystem with public key $n = 1333 = 31 \cdot 43$. The letters of the Latin alphabet A, \ldots, Z are represented by the following 5 bit scheme: A = 00000, B = 00001,..., Z = 11001. Decipher the cryptogram c. Remark: The security requirement to use at most $h = \lfloor \log_2 \lfloor \log_2(n) \rfloor \rfloor$ bits of the Blum-Blum-Shub generator is violated in this example. Instead, 5 bits of the output are used.

Problem 2. (Blum-Blum-Shub generator) The security of the Blum-Blum-Shub generator is based on the difficulty to compute square roots modulo n = pq for two distinct primes p and q with $p, q \equiv 3 \mod 4$.

Design a generator for pseudo-random bits which is based on the hardness of the RSA-problem.

Problem 3. (basic requirements for cryptographic hash functions) Using a block cipher $E_K(x)$ with block length k and key K, a hash function h(m) is provided in the following way:

Append m with zero bits until it is a multiple of k, divide m into n blocks of k bits each.

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c \leftarrow E_{m_0}(m_0)

for i in 1..(n-1) do

d \leftarrow E_{m_0}(m_i)

c \leftarrow c \oplus d

end for

h(m) \leftarrow c
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- a) Does this function fulfill the basic requirements for a cryptographic hash function?
- **b)** Can these requirements be fulfilled by replacing the operation XOR (\oplus) by AND (\odot) ?