Dr. Michael Reyer

## Tutorial 13

Wednesday, February 6, 2019

Problem 1. (Rabin cryptosystem) Consider a Rabin cryptosystem with public key $n^{\prime}=p^{\prime} \cdot q^{\prime}=989$.
a) Check that the given parameters satisfy the Rabin cryptosystem requirements.

Consider a Rabin cryptosystem with another public key $n=p \cdot q=161$ with $p=7$ and $q=23$. In order to be able to identify the correct message, we know that the binary representation of the plaintext $m$ ends with 1111.
b) Decipher the cryptogram $c=116$.
c) Show that the Rabin cryptosystem is vulnerable to chosen-ciphertext attacks.
d) Which vulnerability has the Rabin cryptosystem if $m<\sqrt{n}$ holds? How would you resolve this weakness?

Problem 2. (Merkle signature scheme)
a) What are the four main requirements for cryptographic hash functions?

Let $1<L \in \mathbb{N}$ and $h(m)=m^{2}-1 \bmod L, m \in \mathbb{Z}$, a hash-function.
b) Let $m \in \mathbb{Z}$. Determine an $m \neq m^{\prime} \in \mathbb{Z}$ such that $h(m)=h\left(m^{\prime}\right)$.

Consider the following hash-based signature scheme to sign messages $m \in \mathbb{N}$. Let $\operatorname{bin}(m)$ denote the binary representation of $m \in \mathbb{N}$. Assume that $\operatorname{bin}(m)$ has $n$ bits.

## Key Generation

1) Select $t=n+\left\lfloor\log _{2}(n)\right\rfloor+1$ random numbers $k_{i}$.
2) Compute $v_{i}=h\left(k_{i}\right)$ for all $i=1, \ldots, t$, using a hash function $h: \mathbb{Z} \rightarrow \mathbb{Z}_{L}$ with $L \in \mathbb{N}$.
3) The public key is $\left(v_{1}, v_{2}, \ldots, v_{t}\right)$ and the private key is $\left(k_{1}, k_{2}, \ldots, k_{t}\right)$.

## Signature Generation

1) Compute $\hat{c}$, the binary representation of the number of zeros of $\hat{m}=\operatorname{bin}(m)$.
2) Form the concatenated message $\hat{w}=\hat{m}\left\|\hat{c}=\left(a_{1}, a_{2}, \ldots, a_{n}\right)\right\|\left(a_{n+1}, \ldots, a_{t}\right)$ with bits $a_{i}$, for all $i \leq 1 \leq t$.
3) Determine the positions $i_{1}<i_{2}<\ldots<i_{u}$ in $\hat{w}$, where $a_{i_{j}}=1$, for all $1 \leq j \leq u$.
4) Set $s_{j}=k_{i_{j}}$ for all $1 \leq j \leq u$.
5) The signature for $m$ is $\left(s_{1}, s_{2}, \ldots, s_{u}\right)$.

Solve the following tasks assuming that $\operatorname{bin}(m)$ has $n=5$ bits. The private key is given as $(6,36,27,24,12,3,9,34)$.
c) Describe a verification of the above signature scheme.
d) Sign the decimal message $m=10$.
e) Eve intercepts a sequence of signatures from Alice. Which knowledge is needed by Eve to impersonate Alice and sign arbitrary messages?

Problem 3. (ElGamal signature scheme) Consider an ElGamal signature scheme. The parameters for the signature scheme are given as

$$
p=113, x=66, a=3 .
$$

a) Show that $a$ is a primitive element modulo $p$.
b) Calculate the ElGamal signature for the hash value $h(m)=77$ using the random secret $k=19$.
Hint: $k^{-1} \equiv 59(\bmod p-1)$.
From now on we investigate the verification of the ElGamal signature scheme with general parameters.
Assume initially that message $m$ is signed without using a hash function. Oscar chooses $u, v \in \mathbb{Z}$ with $\operatorname{gcd}(v, p-1)=1$. He computes

$$
\begin{aligned}
& r=a^{u} y^{v} \quad \bmod p, \\
& s=-r v^{-1} \quad \bmod (p-1) .
\end{aligned}
$$

c) Show that $(r, s)$ is a valid signature for the message $m=s u \bmod (p-1)$.

Oscar knows the signature ( $\hat{r}, \hat{s}$ ) for the hash value $h(\hat{m})$. Let $h(\hat{m})$ and $h\left(m^{\prime}\right)$ be invertible modulo ( $p-1$ ). Oscar computes

$$
\begin{aligned}
r^{\prime} & =\hat{r}\left(h\left(m^{\prime}\right) h(\hat{m})^{-1} p-p+1\right) \quad \bmod (p(p-1)), \\
s^{\prime} & =\hat{s} h\left(m^{\prime}\right) h(\hat{m})^{-1} \bmod (p-1) .
\end{aligned}
$$

d) Show that the verification of the signature $\left(r^{\prime}, s^{\prime}\right)$ for $h\left(m^{\prime}\right)$ provides $v_{1}=v_{2}$.
e) Why does this attack still fail if the verification process is correctly applied?

Problem 4. (Elliptic Curve Cryptography) Consider the equation

$$
E_{a}: Y^{2}=X^{3}+a X
$$

You want to uniquely map a message $0<m<\frac{p}{2}$ to a point $\left(x_{m}, y_{m}\right)$ on the elliptic curve $E_{a}\left(\mathbb{F}_{p}\right)$, where $p$ is prime with $p \equiv 3(\bmod 4)$ and $x_{m} \in\{m, p-m\}$.
a) Show that such a point exists. Substantiate your answer.
b) Map the message $m=6$ to a point on the elliptic curve $E_{1}\left(\mathbb{F}_{131}\right)$.

Hint: You may use that $g=2$ is generator of the field $\mathbb{F}_{131}$ and $2^{114} \equiv 91(\bmod 131)$.
Let $p=7$.
c) Determine all $a \in \mathbb{F}_{7}$ such that $E_{a}\left(\mathbb{F}_{7}\right)$ describes an elliptic curve.
d) For which $a \in \mathbb{F}_{7}$ does the point $(3,2)$ lie on $E_{a}\left(\mathbb{F}_{7}\right)$ ?

Let $a=1$.
e) Calculate all points on $E_{1}\left(\mathbb{F}_{7}\right)$
f) What is the order and the trace $t$ of $E_{1}\left(\mathbb{F}_{7}\right)$ ?
g) Prove that $P=(3,3)$ is a generator of the group.

Hint: Use $2 \cdot(3,3)=(1,4)$.

